# On the approximation power of two-layer networks of random ReLUs 

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## Two-layer networks of random ReLUs ("random ReLU networks")



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f \in \operatorname{span}\{\underbrace{x \mapsto \max \left\{0, \mathbf{w}^{(i)} \cdot x-\mathbf{b}^{(i)}\right\}}_{\mathbf{g}^{(i)}}: i \in[r]\}, \quad\left(\left(\mathbf{w}^{(i)}, \mathbf{b}^{(i)}\right)\right)_{i=1}^{r} \sim \mathcal{D}
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## Approximating Lipschitz functions by two-layer networks of random ReLUs

Two-layer networks of random ReLUs:

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\mathcal{F}_{r}:=\operatorname{span}\{\underbrace{x \mapsto \max \left\{0, \mathbf{w}^{(i)} \cdot x-\mathbf{b}^{(i)}\right\}}_{\mathbf{g}^{(i)}}: i \in[r]\}, \quad\left(\left(\mathbf{w}^{(i)}, \mathbf{b}^{(i)}\right)\right)_{i=1}^{r} \sim \mathcal{D},
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where $\mathcal{D}$ is probability distribution for bottom-level parameters $\left(\mathbf{w}^{(i)}, \mathbf{b}^{(i)}\right) \in S^{d-1} \times \mathbb{R}$

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## Question:

What is the minimum width $r$ s.t. $\mathcal{F}_{r}$ can $\varepsilon$-approximate any $L$-Lipschitz functions in $\mathcal{L}^{2}\left([-1,1]^{d}\right)$ (with high probability)?

$$
\operatorname{Pr}\left[\inf _{\hat{f} \in \mathcal{F}_{r}}\left\|\hat{f}-f^{\star}\right\|_{\mathcal{L}^{2}\left([-1,1]^{d}\right)} \leq \varepsilon\right] \geq 0.9 \quad \text { for all } L \text {-Lipschitz } f^{*}:[-1,1]^{d} \rightarrow \mathbb{R}
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\|f\|_{\mathcal{L}^{2}\left([-1,1]^{d}\right)}=\sqrt{\underset{\mathrm{x} \sim \operatorname{Unif}\left([-1,1]^{d}\right)}{\mathbb{E}}\left[f(\mathrm{x})^{2}\right]}
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Our work: upper- and lower-bounds on this minimum width, for all $d, \varepsilon$, and $L$

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1. Approximation capability of neural networks at (or near) random initialization
[Andoni, Panigrahy, Valiant, \& Zhang, '14; Bach, '17; Ji, Telgarsky, \& Xian, '19; Yehudai \& Shamir, '19; ...]
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2. Interplay between dimension $d$ and relative error $\varepsilon / L$


## Our results (informally)

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\geq \exp (\Omega(d)) & \text { if } L / \varepsilon=\Omega(\sqrt{d})
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## Some prior work

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|  | Width | Comments |
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| Maiorov, '99 | $\geq \exp (\Omega(d))$ | $L / \varepsilon \rightarrow \infty$ |
| Yehudai \& Shamir, '19; <br> Kamath, Montasser, \& Srebro, '20 | $\geq \exp (\Omega(d))$ | $L / \varepsilon \geq \operatorname{poly}(d)$ |
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Maiorov's bound (for $H^{1}\left([-1,1]^{d}\right)$ ) applies to networks with arbitrary bottom-level weights, but only holds asymptotically as $L / \varepsilon \rightarrow \infty$

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Hard function of YS and KMS has poly (d) Lipschitz constant

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$\mathcal{L}^{\infty}$ approximation is stronger than $\mathcal{L}^{2}$ approximation

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Upshot: Prior work doesn't reveal the correct minimum width for arbitrary $d$ and $L / \varepsilon$

## Outline for rest of talk

1. Upper- and lower-bounds on the minimum width
2. Proof sketches
3. Some consequences

## Part 1. Upper- and lower-bounds on the minimum width



## Our main results

$$
\operatorname{MinWidth}_{\varepsilon, d, \mathcal{D}}\left(f^{\star}\right):=\min \left\{r \in \mathbb{N}: \operatorname{Pr}\left[\inf _{\hat{f} \in \mathcal{F}_{r}}\left\|\hat{f}-f^{\star}\right\|_{\mathcal{L}^{2}\left([-1,1]^{d}\right)} \leq \varepsilon\right] \geq 0.9\right\}
$$

smallest width $r$ s.t. $\mathcal{F}_{r}$ (with bottom-level weights $\sim \mathcal{D}$ ) $\varepsilon$-approximates $f^{\star}$ with probability $\geq 90 \%$

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Q_{k, d}:=\left|\left\{\alpha \in \mathbb{Z}^{d}:\|\alpha\|_{2} \leq k\right\}\right| \quad \text { number of integer lattice points in radius } k \text { ball in } \mathbb{R}^{d}
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Theorem 1 (upper bound). For any $L, \varepsilon, d$, there exists a parameter distribution $\mathcal{D}$ such that

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Lower-bound, in fact, applies to any target-independent $\mathcal{F}_{r}$ (not just span of random ReLUs)

## Counting integer lattice points in a ball

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Generalized Gauss Circle Problem: As $k \rightarrow \infty$,

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Q_{k, d}=\operatorname{vol}\left(B_{d}\right) \cdot k^{d} \cdot(1+o(1)) \approx \frac{1}{\sqrt{\pi d}}\left(\frac{2 \pi e k^{2}}{d}\right)^{d / 2} \cdot(1+o(1))
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But when $d$ is large compared to $k^{2}$, more favorable bounds are obtained via (simple) combinatorics:

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Theorems $1 \& 2 \Longrightarrow \sup _{L \text {-Lipschitz } f^{\star}} \operatorname{MinWidth}_{\varepsilon, d, \mathcal{D}}\left(f^{\star}\right)= \begin{cases}\operatorname{poly}(d) & \text { if } L / \varepsilon=\Theta(1) \\ \operatorname{poly}(L / \varepsilon) & \text { if } d=\Theta(1) \\ \exp (\Theta(d)) & \text { if } L / \varepsilon=\Theta(\sqrt{d})\end{cases}$

Part 2. Proof sketches


## Proof of upper-bound (sketch)

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1. Get $\varepsilon / 2$-approximation of $L$-Lipschitz $f^{\star}$ using orthonormal basis functions

$$
\sqrt{2} \sin (\pi \alpha \cdot x / 2) \quad \text { and } \quad \sqrt{2} \cos (\pi \alpha \cdot x / 2)
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for $\alpha \in \mathbb{Z}^{d}$ with $\|\alpha\|_{2} \leq 2 L / \varepsilon$

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2. Construct suitable parameter distribution $\mathcal{D}$, so every trigonometric polynomial

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p^{\star} \in \operatorname{span}\left\{\sin (\pi \alpha \cdot x), \cos (\pi \alpha \cdot x): \alpha \in \mathbb{Z}^{d},\|\alpha\|_{2} \leq k\right\}
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with bounded coefficients has

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Basis of "sinusoidal ridge functions" are especially convenient for this step

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We generalize a dimension argument of [Barron, '93]:

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We generalize a dimension argument of [Barron, '93]:

1. If $\varphi_{1}, \ldots, \varphi_{N} \in \mathcal{L}^{2}$ are orthonormal with $N \geq r$, then $\mathcal{F}_{r}$ is $\sqrt{1-\frac{r}{N}}$-far from at least one $\varphi_{i}$

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## Proof of lower-bound (sketch)

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If $\mathcal{D}_{\text {weights }}$ is invariant to coordinate permutations, then the hard-to-approximate function is explicit:

$$
x \mapsto \varepsilon \sin \left(\pi\left(x_{1}+x_{2}+\cdots\right)\right)
$$

## Key lemma

Lemma. Let $H$ be a Hilbert space, and fix orthonormal $\varphi_{1}, \ldots, \varphi_{N} \in H$. Let $\mathbf{W}$ be (possibly random) finite-dimensional subspace of $H$ with $r:=\mathbb{E}[\operatorname{dim}(\mathbf{W})]<\infty$. Then there is some $i \in[N]$ such that

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## Part 3. Some consequences



## Depth separation

- Recent line-of-inquiry on separations between poly-size "shallow" nets and poly-size "deep" nets [Telgarsky, '16; Eldan \& Shamir, '16; Daniely, '17; Safran \& Shamir, '17; Safran, Eldan, \& Shamir, '19; ...]


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Our results $\Rightarrow$ No, for constant $\mathcal{L}^{2}$ approximation error

## Lower-bounds for kernel methods

- Lower-bound applies to all methods that pick $\hat{f}$ from a target-independent subspace of dimension $r$ — including kernel methods based on $r=n$ examples $\left(x^{(1)}, y^{(1)}\right), \ldots,\left(x^{(n)}, y^{(n)}\right)$ :

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- Easy consequence of the key lemma!


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## Thank you!

We gratefully acknowledge support from the NSF (CCF-\{1563155, 1703925, 1740833, 1763970, 1814873\} and IIS-\{1563785, 1838154\}), a Google Faculty Research Award, an Onassis Foundation Scholarship, a Sloan Research Fellowship, and the Simons Collaboration on Algorithms and Geometry.

